

Generalizing BUFIA: Learning Positive and Negative Grammars from Unconventional String Models

Sarah Brogden Payne

sarah.payne@stonybrook.edu

paynesa.github.io



Stony Brook
University



iACS
INSTITUTE FOR ADVANCED
COMPUTATIONAL SCIENCE

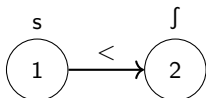
Rutgers Subregular Phonology Workshop

September 29, 2024

- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)



- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations



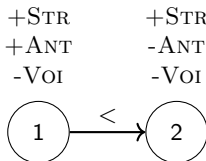
- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations
 - Conversion between negative & positive grammars is straightforward



- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations
 - Conversion between negative & positive grammars is straightforward
- Chandlee et al. (2019): learn with **feature-based representations**



- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations
 - Conversion between negative & positive grammars is straightforward
- Chandee et al. (2019): learn with **feature-based representations**
 - Grammars as collections of **forbidden** combinations



- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations
 - Conversion between negative & positive grammars is straightforward
- Chandlee et al. (2019): learn with **feature-based representations**
 - Grammars as collections of **forbidden** combinations
 - What about collections of **allowed** combinations?



- Previous work: learning formal languages with **symbolic representations** (Heinz, 2010b; Heinz et al., 2012, i.a.)
 - Grammars as collections of **forbidden** or **allowed** combinations
 - Conversion between negative & positive grammars is straightforward
- Chandlee et al. (2019): learn with **feature-based representations**
 - Grammars as collections of **forbidden** combinations
 - What about collections of **allowed** combinations?
- **This talk:** extension of BUFIA to learn grammars as collections of **allowed** or **forbidden** feature-based combinations in a **unified** way



Example: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)

- Subsequences such as [s...s] that agree in \pm ANTERIOR are allowed
- Subsequences such as [s...ʃ] which disagree are banned

✓ [hasxintilawas]

✗ [hasxintilawaʃ]

(Hansson, 2010)

Example: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)

- Subsequences such as $[s\dots s]$ that agree in \pm ANTERIOR are allowed
- Subsequences such as $[s\dots f]$ which disagree are banned

✓ [hasxintilawas]

✗ [hasxintilawaf]

(Hansson, 2010)

Negative Grammar (G^-)

$[+ANT][-ANT] \in G^- \Rightarrow \text{sft} \notin L(G^-)$

Since $[+ANT][-ANT]$ covers $[sf]$

Example: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)

- Subsequences such as $[s\dots s]$ that agree in \pm ANTERIOR are allowed
- Subsequences such as $[s\dots f]$ which disagree are banned

✓ [hasxintilawas]

✗ [hasxintilawaf]

(Hansson, 2010)

Negative Grammar (G^-)

$[+ANT][-ANT] \in G^- \Rightarrow sft \notin L(G^-)$

Since $[+ANT][-ANT]$ covers $[sf]$

Positive Grammar (G^+)

$[+STR][-STR], [-ANT][-ANT] \in G^+ \Rightarrow \text{fft} \in L(G^+)$

Since $[-ANT][-ANT]$ covers $[[ff]$ and $[+STR][-STR]$ covers $[ft]$

Positive Grammars: Tiling

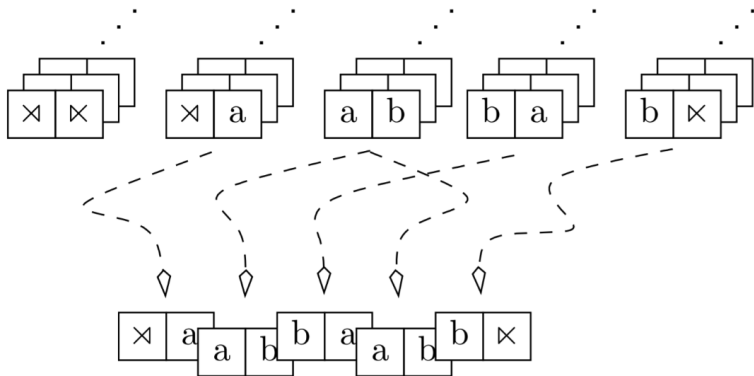


Figure courtesy of Rogers and Heinz (2014)

Why Positive Grammars?

Psycholinguistic Motivation

Computational Motivation

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)

Computational Motivation

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)
- Evidence for positive **syntactic** and **morphological** grammars (Marcus et al., 1992; Yang, 2016; Belth et al., 2021; Li and Schuler, 2023, i.a.)

Computational Motivation

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)
- Evidence for positive **syntactic** and **morphological** grammars (Marcus et al., 1992; Yang, 2016; Belth et al., 2021; Li and Schuler, 2023, i.a.)

Computational Motivation

- **Post-hoc conversion** between positive & negative grammars is straightforward for symbolic models (Heinz, 2010b)

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)
- Evidence for positive **syntactic** and **morphological** grammars (Marcus et al., 1992; Yang, 2016; Belth et al., 2021; Li and Schuler, 2023, i.a.)

Computational Motivation

- **Post-hoc conversion** between positive & negative grammars is straightforward for symbolic models (Heinz, 2010b)
- Grammar polarity has implications for the **learning trajectory**

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)
- Evidence for positive **syntactic** and **morphological** grammars (Marcus et al., 1992; Yang, 2016; Belth et al., 2021; Li and Schuler, 2023, i.a.)

Computational Motivation

- **Post-hoc conversion** between positive & negative grammars is straightforward for symbolic models (Heinz, 2010b)
- Grammar polarity has implications for the **learning trajectory**
- Post-hoc conversion is **exponentially more costly** for models that use features

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments
- Segment-based model: no more than $(s)^k \leq (2^n)^k$ k -factors

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments
- Segment-based model: no more than $(s)^k \leq (2^n)^k$ k -factors
- Under a feature-based model: $(3^n)^k$ possible k -subfactors since each feature can be positive, negative, or unspecified

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments
- Segment-based model: no more than $(s)^k \leq (2^n)^k$ k -factors
- Under a feature-based model: $(3^n)^k$ possible k -subfactors since each feature can be positive, negative, or unspecified

Conversion Process

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments
- Segment-based model: no more than $(s)^k \leq (2^n)^k$ k -factors
- Under a feature-based model: $(3^n)^k$ possible k -subfactors since each feature can be positive, negative, or unspecified

Conversion Process

- For symbolic models, need to simply check whether some k -factor $f \in \Sigma^k$ is in G^-
if $[s\dots j] \in G^-$ then $[s\dots j] \notin G^+$

The Cost of Interdefinability

For symbolic models, negative & positive grammars are straightforwardly interdefinable:

$$G^+ = \Sigma^k \setminus G^- \quad G^- = \Sigma^k \setminus G^+$$

Two complications for feature-based models:

Number of k -Subfactors

- A model with n binary features defines $s \leq 2^n$ segments
- Segment-based model: no more than $(s)^k \leq (2^n)^k$ k -factors
- Under a feature-based model: $(3^n)^k$ possible k -subfactors since each feature can be positive, negative, or unspecified

Conversion Process

- For symbolic models, need to simply check whether some k -factor $f \in \Sigma^k$ is in G^-
if $[s\dots j] \in G^-$ then $[s\dots j] \notin G^+$
- For feature-based models, a k -subfactor f should not be added to G^+ if $f \in G^-$, but also if $(\exists g \in G^-)[f \sqsubseteq g \vee g \sqsubseteq f]$.

Why Positive Grammars?

Psycholinguistic Motivation

- Child may construct **positive phonological** and **phonotactic** grammars (Belth, 2023; Payne, 2023, i.a.)
- Evidence for positive **syntactic** and **morphological** grammars (Marcus et al., 1992; Yang, 2016; Belth et al., 2021; Li and Schuler, 2023, i.a.)

By **fixing k** — the size of the learned substructures — we can straightforwardly adapt the algorithm of Chandlee et al. (2019) to learn the **most general positive and negative grammars** over **feature-based models**

Computational Motivation

- **Post-hoc conversion** between positive & negative grammars is straightforward for symbolic models (Heinz, 2010b)
- Grammar polarity has implications for the **learning trajectory**
- Post-hoc conversion is **exponentially more costly** for models that use features

Table of Contents

- 1 Subfactors and Maxfactors
- 2 Grammars and Their Languages
- 3 The Learning Algorithm
- 4 Example: Samala Sibilant Harmony

Table of Contents

- 1 Subfactors and Maxfactors
- 2 Grammars and Their Languages
- 3 The Learning Algorithm
- 4 Example: Samala Sibilant Harmony

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B
- 2 All relations that hold in A also hold over the corresponding elements in B

Subfactors & Maxfactors

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B
- 2 All relations that hold in A also hold over the corresponding elements in B

An R-structure A is a **maxfactor** of an R-structure B (notated $A \leq B$) if **1 and 2 are satisfied** and:

Subfactors & Maxfactors

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B
- 2 All relations that hold in A also hold over the corresponding elements in B

An R-structure A is a **maxfactor** of an R-structure B (notated $A \leq B$) if **1 and 2 are satisfied** and:

- 3 All relations that hold in B also hold over the corresponding elements in A

Subfactors & Maxfactors

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B
- 2 All relations that hold in A also hold over the corresponding elements in B

An R-structure A is a **maxfactor** of an R-structure B (notated $A \leq B$) if **1 and 2 are satisfied** and:

- 3 All relations that hold in B also hold over the corresponding elements in A

**Subfactor:
Unidirectional**

Subfactors & Maxfactors

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if:

- 1 There is a mapping between D^A and some subset of D^B
- 2 All relations that hold in A also hold over the corresponding elements in B

An R-structure A is a **maxfactor** of an R-structure B (notated $A \leq B$) if **1 and 2 are satisfied** and:

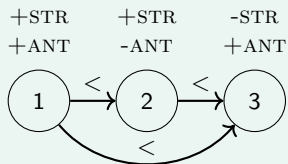
- 3 All relations that hold in B also hold over the corresponding elements in A

**Subfactor:
Unidirectional**

**Maxfactor:
Bidirectional**

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

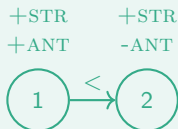
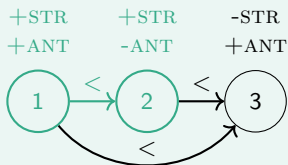


Subfactor:
Unidirectional

Maxfactor:
Bidirectional

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

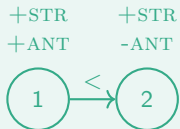
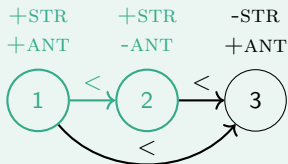


Subfactor:
Unidirectional

Maxfactor:
Bidirectional

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

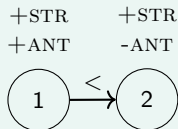
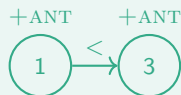
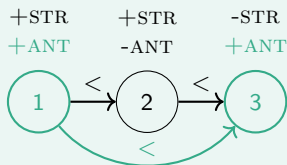


✓ **Subfactor:**
Unidirectional

✓ **Maxfactor:**
Bidirectional

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

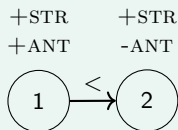
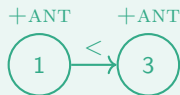
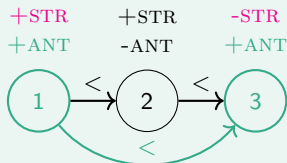


Subfactor:
Unidirectional

Maxfactor:
Bidirectional

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

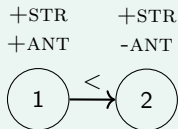
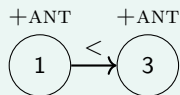
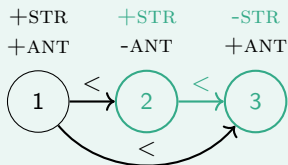


✓ **Subfactor:**
Unidirectional

~~X~~ **Maxfactor:**
Bidirectional

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)

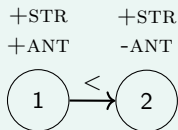
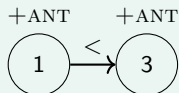
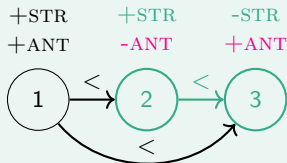


**Subfactor:
Unidirectional**

**Maxfactor:
Bidirectional**

Subfactors vs. Maxfactors

Examples (Samala Sibilant Harmony)



✓ **Subfactor:**
Unidirectional

✗ **Maxfactor:**
Bidirectional

Definition: k -Subfactors

If $A \sqsubseteq B$ and $|A| = k$, then A is a **k -subfactor** of B

Definition: k -Maxfactors

If $A \leq B$ and $|A| = k$, then A is a **k -maxfactor** of B

Definition: k -Subfactors

If $A \sqsubseteq B$ and $|A| = k$, then A is a k -**subfactor** of B

Let the set of k -subfactors of an R -structure B be given by:

$$\text{SFAC}_k(B) := \{A \mid A \sqsubseteq B, |A| = k\}$$

Definition: k -Maxfactors

If $A \leq B$ and $|A| = k$, then A is a k -**maxfactor** of B

Let the set of k -maxfactors of B be given by:

$$\text{MFAC}_k(B) := \{A \mid A \leq B, |A| = k\}$$

Table of Contents

- 1 Subfactors and Maxfactors
- 2 Grammars and Their Languages**
- 3 The Learning Algorithm
- 4 Example: Samala Sibilant Harmony

Positive vs. Negative Grammars

Grammar G = finite set of k -subfactors

Language defined by G depends on its **interpretation**:

Positive vs. Negative Grammars

Grammar G = finite set of k -subfactors

Language defined by G depends on its **interpretation**:

Negative Grammar (G^-)

Elements of G^- are forbidden,
and strings in $L(G^-)$ contain
no forbidden subfactors

Positive vs. Negative Grammars

Grammar G = finite set of k -subfactors

Language defined by G depends on its **interpretation**:

Negative Grammar (G^-)

Elements of G^- are forbidden,
and strings in $L(G^-)$ contain
no forbidden subfactors

Positive Grammar (G^+)

Elements of G^+ are
permissible, and strings in
 $L(G^+)$ are those which are
tiled by these elements

Languages of Positive vs. Negative Grammars

Negative Grammar

The language $L(G^-)$ of G^- is given by:

$$L(G^-) = \{w \in \Sigma^* \mid (\forall S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^- = \emptyset]\}$$

or equivalently by:

$$L(G^-) = \{w \in \Sigma^* \mid (\nexists S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^- \neq \emptyset]\}$$

Positive Grammar

The language $L(G^+)$ of G^+ is given by:

$$L(G^+) = \{w \in \Sigma^* \mid (\forall S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^+ \neq \emptyset]\}$$

or equivalently by:

$$L(G^+) = \{w \in \Sigma^* \mid (\nexists S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^+ = \emptyset]\}$$

Languages of Positive vs. Negative Grammars

Negative Grammar

The language $L(G^-)$ of G^- is given by:

$$L(G^-) = \{w \in \Sigma^* \mid (\forall S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^- = \emptyset]\}$$

or equivalently by:

$$L(G^-) = \{w \in \Sigma^* \mid (\nexists S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^- \neq \emptyset]\}$$

$\in G$

Positive Grammar

The language $L(G^+)$ of G^+ is given by:

$$L(G^+) = \{w \in \Sigma^* \mid (\forall S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^+ \neq \emptyset]\}$$

or equivalently by:

$$L(G^+) = \{w \in \Sigma^* \mid (\nexists S \in \text{MFAC}_k(M, w)) \\ [\text{SFAC}_k(S) \cap G^+ = \emptyset]\}$$

$\notin G$

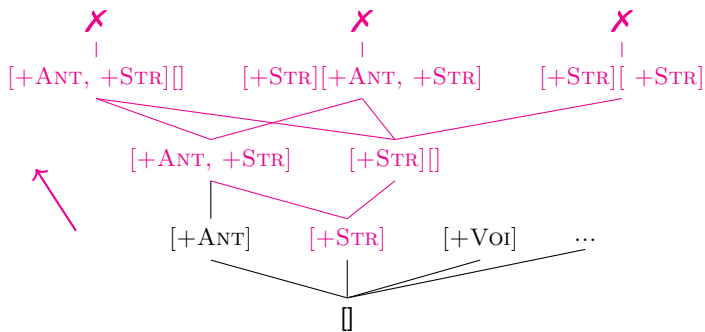
\forall	Positive Grammar	Negative Grammar
\exists	Negative Grammar	Positive Grammar

Table of Contents

- 1 Subfactors and Maxfactors
- 2 Grammars and Their Languages
- 3 The Learning Algorithm**
- 4 Example: Samala Sibilant Harmony

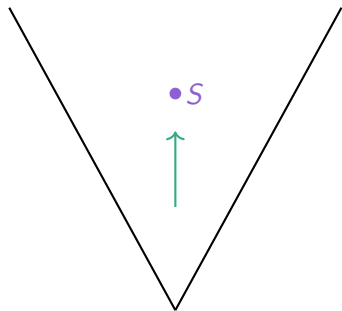
Previous Work

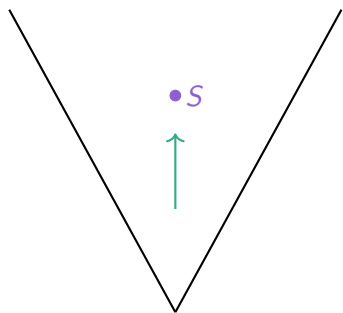
Crucial insight of **Chandlee et al. (2019)**: **grammatical entailment**



For a **negative grammar**:

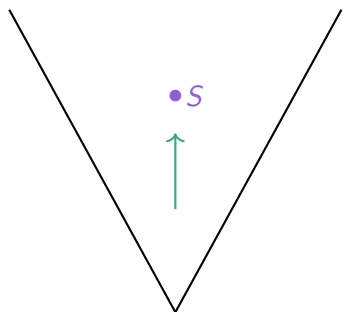
- Add S to G^- if $S \not\sqsubseteq x$ for any $x \in D$





For a **negative grammar**:

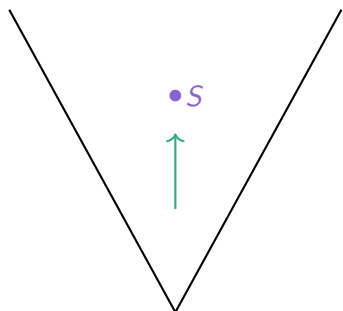
- Add S to G^- if $S \not\sqsubseteq x$ for any $x \in D$
- **Equivalent:** Given the set of all maxfactors that are superfactors of S , **none are attested** in D



For a **negative grammar**:

- Add S to G^- if $S \not\sqsubseteq x$ for any $x \in D$
- **Equivalent:** Given the set of all maxfactors that are superfactors of S , **none are attested** in D

	$\in D$	$\notin D$
\forall	Positive Grammar	Negative Grammar
\exists	Negative Grammar	Positive Grammar



For a **negative grammar**:

- Add S to G^- if $S \not\sqsubseteq x$ for any $x \in D$
- **Equivalent:** Given the set of all maxfactors that are superfactors of S , **none are attested** in D

	$\in D$	$\notin D$
\forall	Positive Grammar	Negative Grammar
\exists	Negative Grammar	Positive Grammar

For a **positive grammar**:

- Given the set of all maxfactors that are superfactors of S , **all are attested** in D

Extensions of a Subfactor

The extensions of a subfactor S are defined as follows:

$$\text{EXT}_k(S) = \{A \in \text{SFAC}_k(M, \Sigma^*) \mid S \sqsubseteq A \wedge (\nexists A') [|A'| = k \wedge A \sqsubseteq A']\} \quad (1)$$

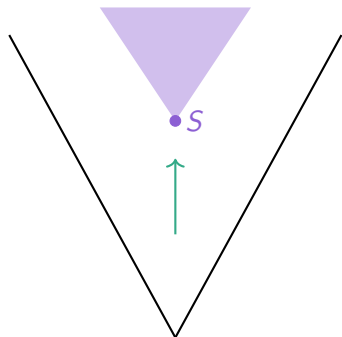
Intuition: extensions of S are all k -maxfactors that are superfactors of S .

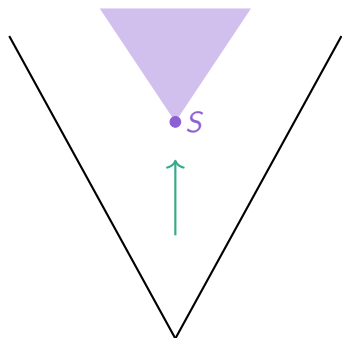
Examples (Extension of $[+ANT]$)

If $S = [+ANT]$ and the only features available are $\pm ANT$, $\pm VOI$ and $\pm STR$:

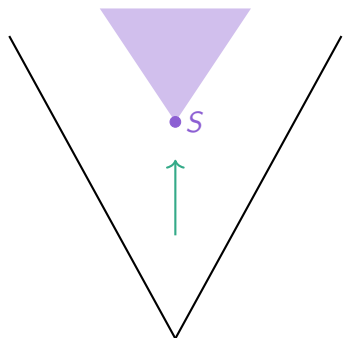
$$\text{EXT}_k(S) = \{ [+ANT, -STR, +VOI], [+ANT, +STR, +VOI], [+ANT, -STR, -VOI], [+ANT, +STR, -VOI] \}$$

- **Bottom-up** traversal



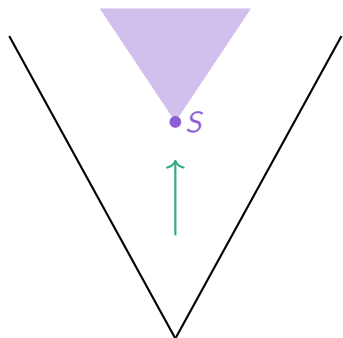


- **Bottom-up** traversal
- For a given subfactor S , check whether:



- **Bottom-up** traversal
- For a given subfactor S , check whether:
 - G is **negative** and

$$(\forall s' \in \text{EXT}_k(S))[\nexists x \in D, s' \leq x]$$



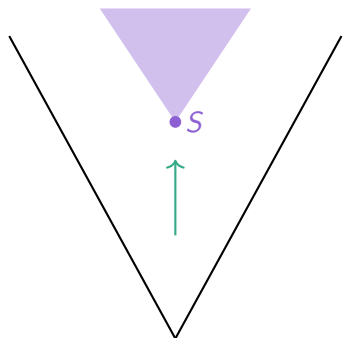
- **Bottom-up** traversal
- For a given subfactor S , check whether:

- G is **negative** and

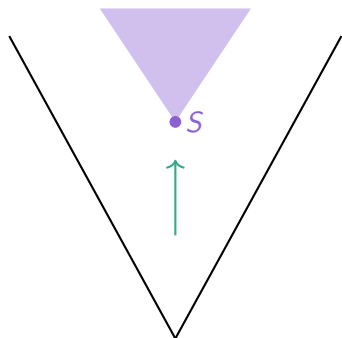
$$(\forall s' \in \text{EXT}_k(S))[\nexists x \in D, s' \leq x]$$

- G is **positive** and

$$(\forall s' \in \text{EXT}_k(S))[\exists x \in D, s' \leq x]$$



- **Bottom-up** traversal
- For a given subfactor S , check whether:
 - G is **negative** and
$$(\forall s' \in \text{EXT}_k(S))[\nexists x \in D, s' \leq x]$$
 - G is **positive** and
$$(\forall s' \in \text{EXT}_k(S))[\exists x \in D, s' \leq x]$$
- If either condition is met:
 - **Add S to G !**



- **Bottom-up** traversal
- For a given subfactor S , check whether:
 - G is **negative** and
$$(\forall s' \in \text{EXT}_k(S))[\nexists x \in D, s' \leq x]$$
 - G is **positive** and
$$(\forall s' \in \text{EXT}_k(S))[\exists x \in D, s' \leq x]$$
- If either condition is met:
 - **Add S to G !**
- Otherwise:
 - Add the **least superfactors** of S to the queue to be considered next

Next Superfactor

We extract the **least superfactors** of S — those that differ minimally from S — by calling $\text{NextSupFact}(s)$ where $\text{NextSupFact}()$ is defined as follows:

$$\text{NextSupFact}(S) = \{A \in \text{SFAC}_k(M, \Sigma^*) \mid S \subseteq A \wedge (\nexists A') [S \subseteq A' \subseteq A]\} \quad (2)$$

Intuition: $\text{NextSupFact}()$ returns the superfactors of S that differ minimally from S .

Examples (Next superfactors of $[+ANT]$)

If $S = [+ANT]$ and the only features available are $\pm ANT$, $\pm VOI$ and $\pm STR$:

$$\text{NextSupFact}(S) = \{ [+ANT, -STR], [+ANT, +STR], [+ANT, -VOI], [+ANT, +VOI] \}$$

Fix Σ , model M , positive integer k , and polarity p . For any language $L \in \mathcal{L}^p(M, k)$ and for any finite sample $D \subseteq L$, return a grammar G^p such that:

Fix Σ , model M , positive integer k , and polarity p . For any language $L \in \mathcal{L}^p(M, k)$ and for any finite sample $D \subseteq L$, return a grammar G^p such that:

- 1 G^p is **consistent**, that is, $D \subseteq L(G^p)$.

Fix Σ , model M , positive integer k , and polarity p . For any language $L \in \mathcal{L}^p(M, k)$ and for any finite sample $D \subseteq L$, return a grammar G^p such that:

- 1 G^p is **consistent**, that is, $D \subseteq L(G^p)$.
- 2 $L(G^p)$ is a **smallest language** in $\mathcal{L}^p(M, k)$ which covers D , so that for all $L \in \mathcal{L}^p(M, k)$ where $D \subseteq L$, we have $L(G^p) \subseteq L$.

Fix Σ , model M , positive integer k , and polarity p . For any language $L \in \mathcal{L}^p(M, k)$ and for any finite sample $D \subseteq L$, return a grammar G^p such that:

- 1 G^p is **consistent**, that is, $D \subseteq L(G^p)$.
- 2 $L(G^p)$ is a **smallest language** in $\mathcal{L}^p(M, k)$ which covers D , so that for all $L \in \mathcal{L}^p(M, k)$ where $D \subseteq L$, we have $L(G^p) \subseteq L$.
- 3 G^p includes R-structures S that are restrictions of R-structures S' in other grammars G' that also satisfy (1) and (2). That is, for all G' satisfying (1) and (2) and for all $S' \in G'$, there exists some $S \in G^p$ such that $S \subseteq S'$.

Table of Contents

- 1 Subfactors and Maxfactors
- 2 Grammars and Their Languages
- 3 The Learning Algorithm
- 4 Example: Samala Sibilant Harmony**

Examples (Samala Sibilant Harmony)

Simplifying Assumptions:

Examples (Samala Sibilant Harmony)

Simplifying Assumptions:

- Two features: \pm ANT (relevant) and \pm VOI (irrelevant)

Examples (Samala Sibilant Harmony)

Simplifying Assumptions:

- Two features: \pm ANT (relevant) and \pm VOI (irrelevant)
- $k = 2$

Examples (Samala Sibilant Harmony)

Simplifying Assumptions:

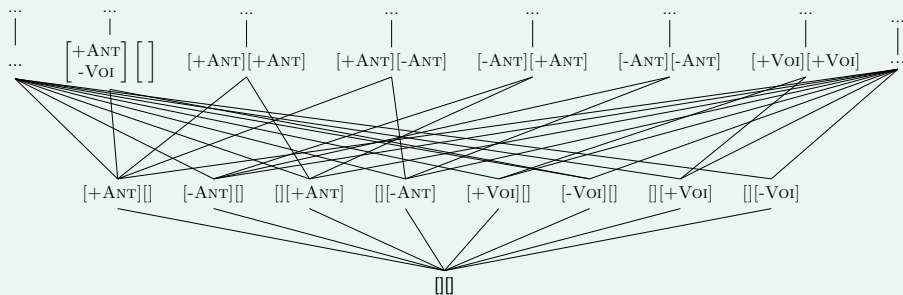
- Two features: \pm ANT (**relevant**) and \pm VOI (**irrelevant**)
- $k = 2$
- All licit subsequences are attested (cf. Heinz, 2010a)

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)

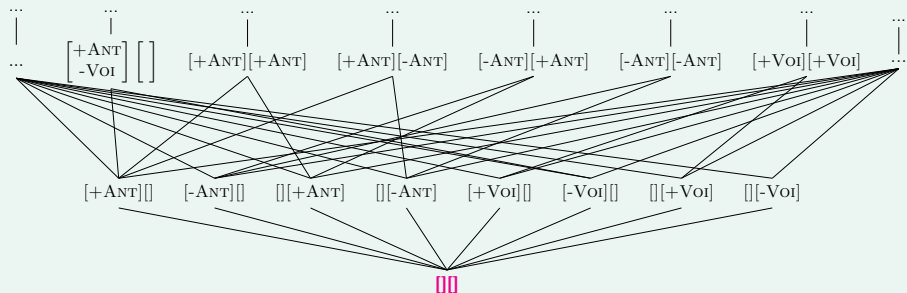
Simplifying Assumptions:

- Two features: $\pm\text{ANT}$ (relevant) and $\pm\text{VOI}$ (irrelevant)
- $k = 2$
- All licit subsequences are attested (cf. Heinz, 2010a)



Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)

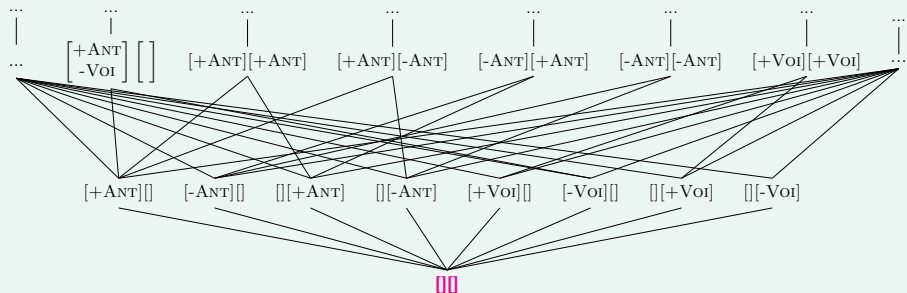


Negative Grammar

Positive Grammar

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



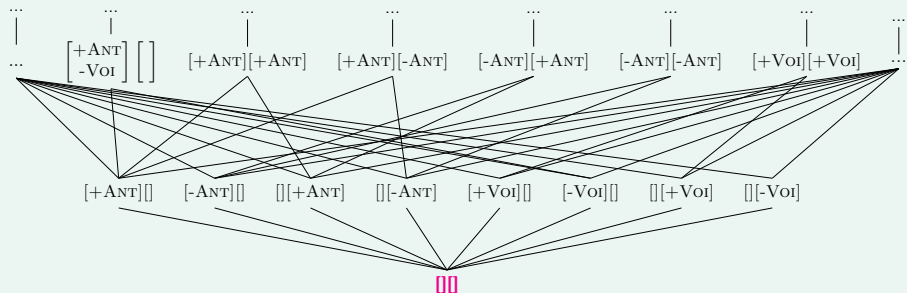
Negative Grammar

Is there any element in $EXT_k([])$ which is a 2-maxfactor of some $x \in D$?

Positive Grammar

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

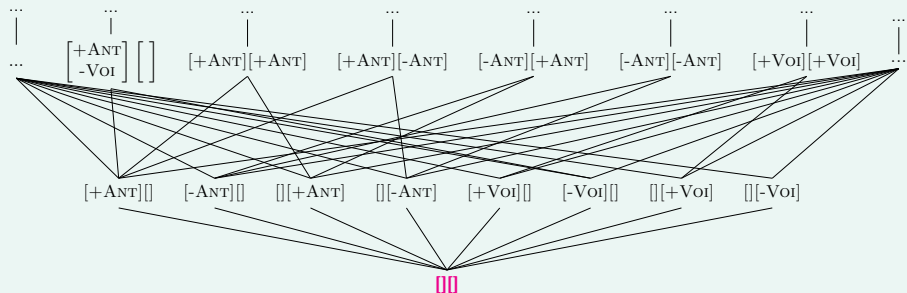
Is there any element in $\text{EXT}_k([\])$ which **is** a 2-maxfactor of some $x \in D$?

Positive Grammar

Is there any element in $\text{EXT}_k([\])$ which **is not** a 2-maxfactor of some $x \in D$?

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

Is there any element in $\text{EXT}_k(\emptyset\emptyset)$ which **is** a 2-maxfactor of some $x \in D$?

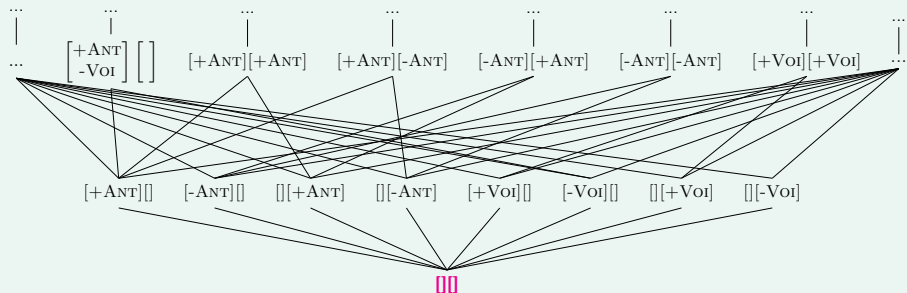
e.g. $[+VOI, +ANT][+VOI, +ANT] \in \text{EXT}_k(\emptyset\emptyset)$, but $[z\dots z]$ is **licit and attested**.

Positive Grammar

Is there any element in $\text{EXT}_k(\emptyset\emptyset)$ which **is not** a 2-maxfactor of some $x \in D$?

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

Is there any element in $\text{EXT}_k(\square\square)$ which **is** a 2-maxfactor of some $x \in D$?

e.g. $[+VOI, +ANT][+VOI, +ANT] \in \text{EXT}_k(\square\square)$, but $[z\dots z]$ is **licit and attested**.

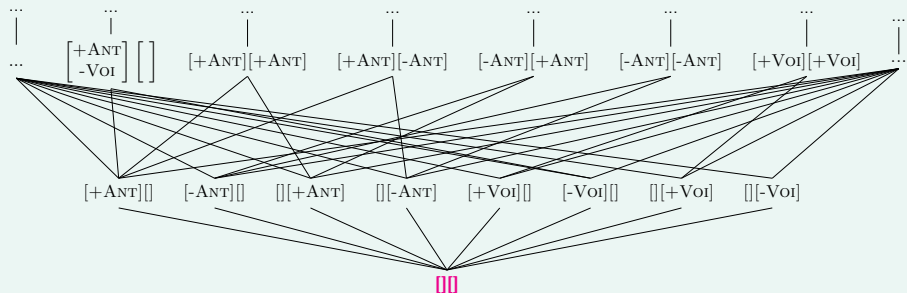
Positive Grammar

Is there any element in $\text{EXT}_k(\square\square)$ which **is not** a 2-maxfactor of some $x \in D$?

e.g. $[+VOI, +ANT][+VOI, -ANT] \in \text{EXT}_k(\square\square)$, but $[z\dots z]$ is **illicit and unattested**.

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

Is there any element in $\text{EXT}_k([\])$ which **is** a 2-maxfactor of some $x \in D$?

e.g. $[+VOI, +ANT][+VOI, +ANT] \in \text{EXT}_k([\])$, but $[z\dots z]$ is **licit and attested**.

Positive Grammar

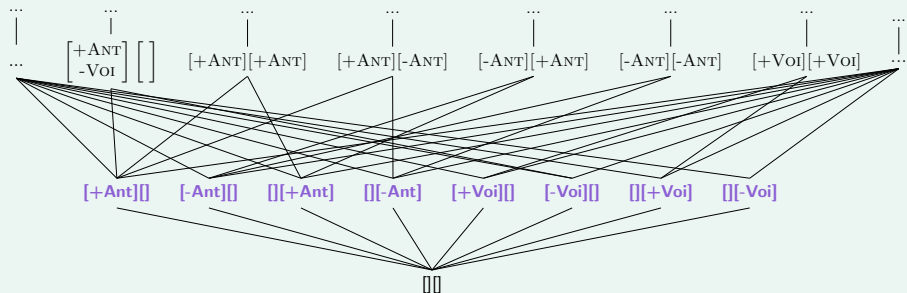
Is there any element in $\text{EXT}_k([\])$ which **is not** a 2-maxfactor of some $x \in D$?

e.g. $[+VOI, +ANT][+VOI, -ANT] \in \text{EXT}_k([\])$, but $[z\dots z]$ is **illicit and unattested**.

Keep searching!

Applying the Algorithm: Samala Sibilant Harmony

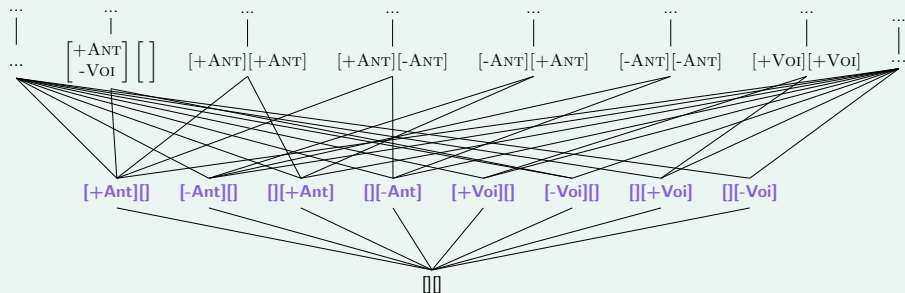
Examples (Samala Sibilant Harmony)



- Extract the **least superfactors** of $[]$ and consider each of them

Applying the Algorithm: Samala Sibilant Harmony

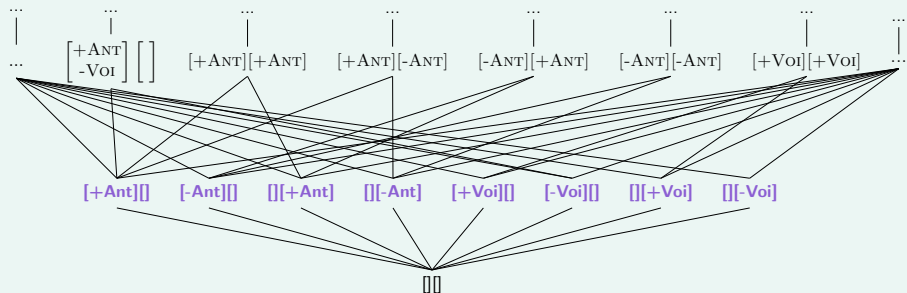
Examples (Samala Sibilant Harmony)



- Extract the **least superfactors** of $[] []$ and consider each of them
- Still not specified enough:

Applying the Algorithm: Samala Sibilant Harmony

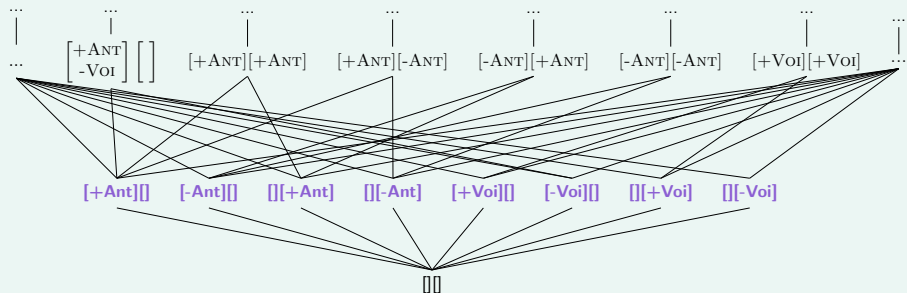
Examples (Samala Sibilant Harmony)



- Extract the **least superfactors** of $[]$ and consider each of them
- Still not specified enough:
 - Any subfactor with $\pm ANT$ specified in one position has licit and illicit maxfactors in its extension

Applying the Algorithm: Samala Sibilant Harmony

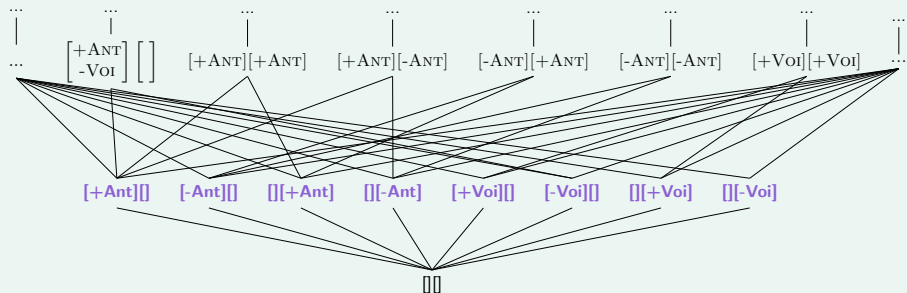
Examples (Samala Sibilant Harmony)



- Extract the **least superfactors** of $[\]$ and consider each of them
- Still not specified enough:
 - Any subfactor with $\pm ANT$ specified in one position has licit and illicit maxfactors in its extension
 $[+ANT][\] \sqsubseteq [+ANT][+ANT]$ but $[+ANT][\] \sqsubseteq [+ANT][-ANT]$

Applying the Algorithm: Samala Sibilant Harmony

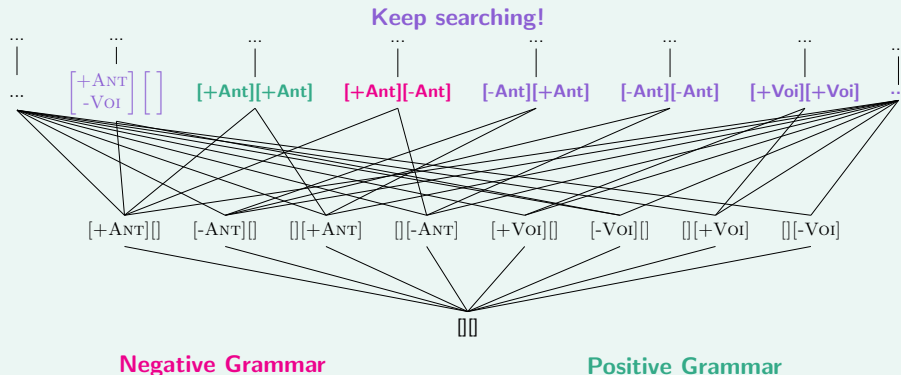
Examples (Samala Sibilant Harmony)



- Extract the **least superfactors** of $[] []$ and consider each of them
- Still not specified enough:
 - Any subfactor with $\pm \text{ANT}$ specified in one position has licit and illicit maxfactors in its extension
 $[+\text{ANT}] [] \sqsubseteq [+\text{ANT}] [+ \text{ANT}]$ but $[+\text{ANT}] [] \sqsubseteq [+\text{ANT}] [- \text{ANT}]$
 - $\pm \text{VOI}$ has no bearing on licitness

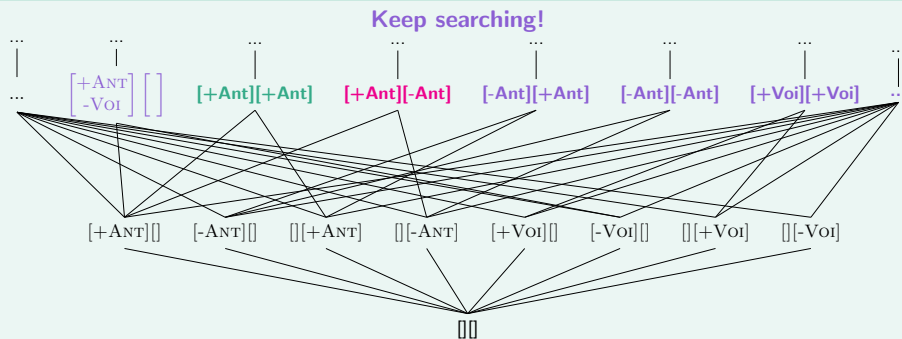
Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

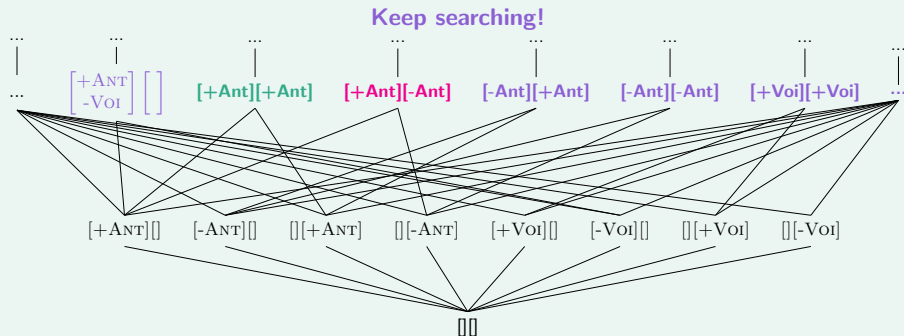
Is there any element in $EXT_k([+ANT][-ANT])$ which is a 2-maxfactor of some $x \in D$? **No!**

Positive Grammar

Is there any element in $EXT_k([+ANT][+ANT])$ which is not a 2-maxfactor of some $x \in D$? **No!**

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

Is there any element in $EXT_k([+ANT][-ANT])$ which **is** a 2-maxfactor of some $x \in D$? **No!**

$[+ANT][-ANT]$ is added to G^-

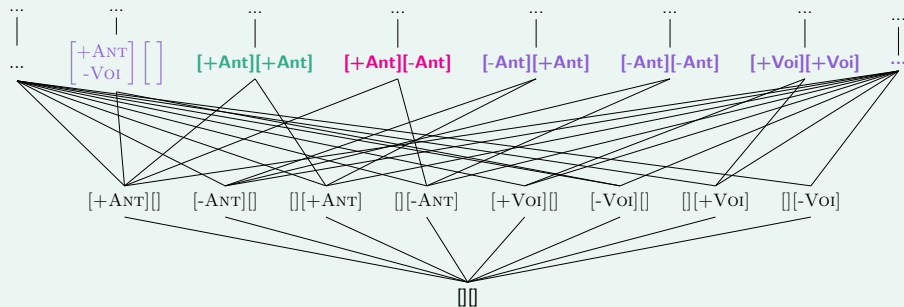
Positive Grammar

Is there any element in $EXT_k([+ANT][+ANT])$ which **is not** a 2-maxfactor of some $x \in D$? **No!**

$[+ANT][+ANT]$ is added to G^+

Applying the Algorithm: Samala Sibilant Harmony

Examples (Samala Sibilant Harmony)



Negative Grammar

We may later reach $[+ANT][-ANT, +VOI]$ but we won't consider it.

$[+ANT][-ANT]$ being banned entails $[+ANT][-ANT, +VOI]$ being banned

Positive Grammar

We may later reach $[+ANT][+ANT, +VOI]$ but we won't consider it.

$[+ANT][+ANT]$ being allowed entails $[+ANT][+ANT, +VOI]$ being allowed

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way
- Enriched representations of feature-based string models allow us to provably find the **most general subfactors**

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way
- Enriched representations of feature-based string models allow us to provably find the **most general subfactors**
- Implications of **grammar polarity**:

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way
- Enriched representations of feature-based string models allow us to provably find the **most general subfactors**
- Implications of **grammar polarity**:
 - As G^+ grows, $L(G^+)$ **grows**

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way
- Enriched representations of feature-based string models allow us to provably find the **most general subfactors**
- Implications of **grammar polarity**:
 - As G^+ grows, $L(G^+)$ **grows**
 - As G^- grows, $L(G^-)$ **shrinks**

- If we fix the size k of subfactors in the grammar, BUFIA can be adapted to learn **positive** and **negative** grammars in a **unified** way
- Enriched representations of feature-based string models allow us to provably find the **most general subfactors**
- Implications of **grammar polarity**:
 - As G^+ grows, $L(G^+)$ **grows**
 - As G^- grows, $L(G^-)$ **shrinks**
 - Initially, G^+ allows **nothing**, while G^- allows **everything**

- **Implementing** the generalized algorithm and applying it to corpus data

- **Implementing** the generalized algorithm and applying it to corpus data
- Approaches to dealing with **noise** and **sparsity** in the data

- **Implementing** the generalized algorithm and applying it to corpus data
- Approaches to dealing with **noise** and **sparsity** in the data
- Learning mixed grammars containing both **banned** and **allowed** subfactors

- **Implementing** the generalized algorithm and applying it to corpus data
- Approaches to dealing with **noise** and **sparsity** in the data
- Learning mixed grammars containing both **banned** and **allowed** subfactors
- Comparison of **learning trajectories** of **positive** & **negative** grammars

- **Implementing** the generalized algorithm and applying it to corpus data
- Approaches to dealing with **noise** and **sparsity** in the data
- Learning mixed grammars containing both **banned** and **allowed** subfactors
- Comparison of **learning trajectories** of **positive** & **negative** grammars
 - Within a single search of the hypothesis space

- **Implementing** the generalized algorithm and applying it to corpus data
- Approaches to dealing with **noise** and **sparsity** in the data
- Learning mixed grammars containing both **banned** and **allowed** subfactors
- Comparison of **learning trajectories** of **positive** & **negative** grammars
 - Within a single search of the hypothesis space
 - When applied to incrementally larger data sets as a proxy for incremental learning

Thank you!

I am grateful to Jeff Heinz, Thomas Graf, Jon Rawski, Logan Swanson, and the SCiL reviewers for discussion.

This work was supported by the Institute for Advanced Computational Science Graduate Research Fellowship and the National Science Foundation Graduate Research Fellowship Program.



Stony Brook
University



iACS
INSTITUTE FOR ADVANCED
COMPUTATIONAL SCIENCE



- Caleb Belth. 2023. *Towards an Algorithmic Account of Phonological Rules and Representations*. Ph.D. thesis, University of Michigan.
- Caleb Belth, Sarah Payne, Deniz Beser, Jordan Kodner, and Charles Yang. 2021. The greedy and recursive search for morphological productivity. *Proceedings of the 43rd annual meeting of the Cognitive Science Society*, 43:2869–2875.
- J Richard Büchi. 1960. Weak second-order arithmetic and finite automata. *Mathematical Logic Quarterly*, 6(1-6).
- Jane Chandlee, Remi Eyraud, Jeffrey Heinz, Adam Jardine, and Jonathan Rawski. 2019. Learning with partially ordered representations. In *Proceedings of the 16th Meeting on the Mathematics of Language*, pages 91–101, Toronto, Canada. Association for Computational Linguistics.
- Gunnar Ólafur Hansson. 2010. *Consonant harmony: Long-distance interactions in phonology*, volume 145. University of California Press.
- Jeffrey Heinz. 2010a. Learning long-distance phonotactics. *Linguistic Inquiry*, 41(4):623–661.
- Jeffrey Heinz. 2010b. String extension learning. In *Proceedings of the 48th Annual Meeting of the Association for Computational Linguistics*, pages 897–906, Uppsala, Sweden. Association for Computational Linguistics.

- Jeffrey Heinz, Anna Kasprzik, and Timo Kötzing. 2012. Learning in the limit with lattice-structured hypothesis spaces. *Theoretical Computer Science*, 457:111–127.
- Daoxin Li and Kathryn D Schuler. 2023. Acquiring recursive structures through distributional learning. *Language Acquisition*, pages 1–14.
- Gary F Marcus, Steven Pinker, Michael Ullman, Michelle Hollander, T John Rosen, Fei Xu, and Harald Clahsen. 1992. Overregularization in language acquisition. *Monographs of the society for research in child development*, pages i–178.
- Robert McNaughton and Seymour A Papert. 1971. *Counter-Free Automata (MIT research monograph no. 65)*. The MIT Press.
- Sarah Payne. 2023. Marginal sequences are licit but unproductive. Poster presented at the 2023 Annual Meeting of Phonology.
- James Rogers and Jeffrey Heinz. 2014. Model theoretic phonology. In *Workshop slides in the 26th European Summer School in Logic, Language and Information*.
- James Rogers, Jeffrey Heinz, Margaret Fero, Jeremy Hurst, Dakotah Lambert, and Sean Wibel. 2013. Cognitive and sub-regular complexity. In *Formal Grammar: 17th and 18th International Conferences, Revised Selected Papers*, pages 90–108. Springer.

References III

- Kristina Strother-Garcia, Jeffrey Heinz, and Hyun Jin Hwangbo. 2016. Using model theory for grammatical inference: a case study from phonology. In *Proceedings of The 13th International Conference on Grammatical Inference*, pages 66–78.
- Mai H Vu, Ashkan Zehfroosh, Kristina Strother-Garcia, Michael Sebok, Jeffrey Heinz, and Herbert G Tanner. 2018. Statistical relational learning with unconventional string models. *Frontiers in Robotics and AI*, 5:76.
- Charles Yang. 2016. *The price of linguistic productivity: How children learn to break the rules of language*. MIT press.

Lemma 1: Maxfactor-Subfactor Containment

Let k be some positive integer and let M be some model of Σ^* . For any $w \in \Sigma^*$ and for any $F \in \text{SFAC}_k(M, w)$, we have that:

$$[\exists G \in \text{MFAC}_k(M, w)](F \sqsubseteq G)$$

Lemma 2: Union of Subfactors of Maxfactors

Let k be some positive integer and let M be some model of Σ^* . For any $w \in \Sigma^*$, we have that:

$$\bigcup_{S \in \text{MFAC}_k(M, w)} \text{SFAC}_k(S) = \text{SFAC}_k(M, w)$$

Model Signature: a set of relations $R = \{R_1, R_2, \dots, R_n\}$

- Each R_i is an m_i -ary relation

Model Signature: a set of relations $R = \{R_1, R_2, \dots, R_n\}$

- Each R_i is an m_i -ary relation

R-Structure: a tuple of elements $S = \langle D; R_1, R_2, \dots, R_n \rangle$

- D , **the domain**, is a finite set of elements
- Each R_i is a subset of D^{m_i}

Model Signature: a set of relations $R = \{R_1, R_2, \dots, R_n\}$

- Each R_i is an m_i -ary relation

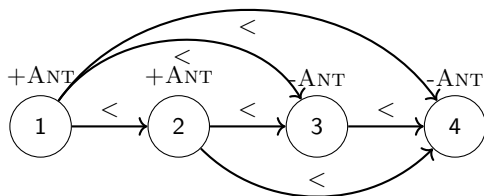
R-Structure: a tuple of elements $S = \langle D; R_1, R_2, \dots, R_n \rangle$

- D , **the domain**, is a finite set of elements
- Each R_i is a subset of D^{m_i}

Size $|S|$ of an R-structure = cardinality of its domain

Precedence Model: $M^<(w) := \langle D^w; <, [R_\sigma^w]_{\sigma \in \Sigma} \rangle$

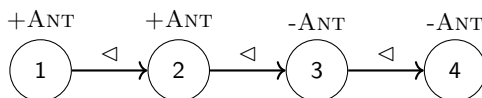
- $D^w = \{1, \dots, |w|\}$ is the **domain** of positions in w
- $< := \{(i, j) \in D^w \times D^w \mid i < j\}$ is the **general precedence** relation



(Büchi, 1960; McNaughton and Papert, 1971; Rogers et al., 2013)

Successor Model: $M^{\triangleleft}(w) := \langle D^w; \triangleleft, [R_{\sigma}^w]_{\sigma \in \Sigma} \rangle$

- $D^w = \{1, \dots, |w|\}$ is the **domain** of positions in w
- $\triangleleft := \{(i, i+1) \in D^w \times D^w\}$ is the **successor** relation



(Büchi, 1960; McNaughton and Papert, 1971; Rogers et al., 2013)

Restrictions

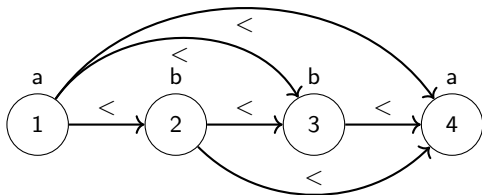
Definition: Restriction

An R-structure A is a **restriction** of an R-structure B if $D^A \subseteq D^B$ and for each m -ary relation R_i in the model signature:

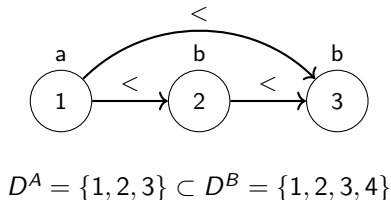
$$R_i^A = \{(x_1, \dots, x_m) \in R_i^B \mid x_1, \dots, x_m \in D^A\} \quad (3)$$

Intuition: identify a subset D^A of the domain of B and retain only those relations in B whose elements are wholly within D^A

R-structure B



R-structure A



Definition: Subfactor

An R-structure A is a **subfactor** of an R-structure B (notated $A \sqsubseteq B$) if there exists a restriction B' of B and a bijection h such that for all $R_i \in R$, if $R_i(x_1, \dots, x_m)$ holds in A , then $R_i(h(x_1), \dots, h(x_m))$ holds in B' .

Intuition: A is a subfactor of B if there is a mapping between D^A and some subset of D^B and all relations that hold in A also hold over the corresponding elements in B

Definition: Maxfactor

An R-structure A is a **maxfactor** of an R-structure B (notated $A \leq B$) iff $A \sqsubseteq B$ and for each m -ary relation R_i , whenever $R_i(x_1, \dots, x_m)$ holds in B , $R_i(h^{-1}(x_1), \dots, h^{-1}(x_m))$ holds in A .

Intuition: A is a maxfactor of B if $A \sqsubseteq B$ and all relations that hold in B also hold over the corresponding elements in A

Connectedness

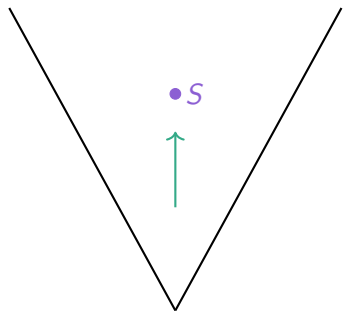
An R-structure $S = \langle D; R_1, R_2, \dots, R_n \rangle$ is connected iff $(\forall x, y \in D)[(x, y) \in C^*]$, where C^* is defined as the symmetric transitive closure of:

$$C = \{(x, y) \in D \times D \mid \\ \exists i \in \{1 \dots n\}, \exists (x_1 \dots x_m) \in R_i \\ \exists s, t \in \{1 \dots m\}, x = x_s, y = x_t\}$$

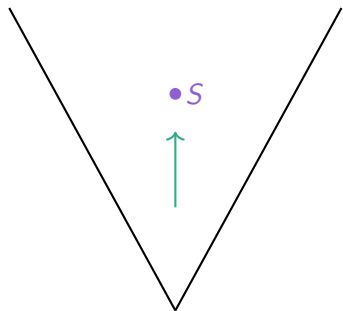
Intuition: domain elements x and y of S belong to C if they belong to some non-unary relation R_i in S

Examples (Disconnected R-Structure)

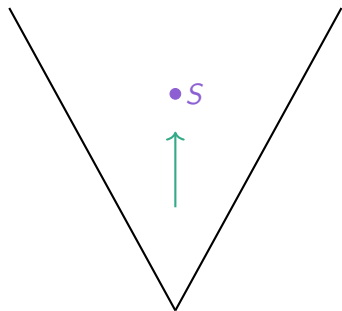




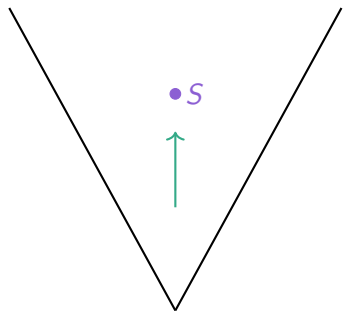
- **Bottom-up** traversal



- **Bottom-up** traversal
- For a given subfactor S , check whether $S \sqsubseteq x$ for any x in the data D



- **Bottom-up** traversal
- For a given subfactor S , check whether $S \sqsubseteq x$ for any x in the data D
 - If not, **posit a constraint:** $S \in G^-$

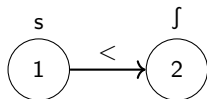


- **Bottom-up** traversal
- For a given subfactor S , check whether $S \sqsubseteq x$ for any x in the data D
 - If not, **posit a constraint**: $S \in G^-$
 - If so, **cannot posit a constraint**
Add the **least superfactors** of S to the queue to be considered next

Conventional vs. Unconventional String Models

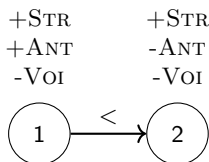
Conventional String Models

- **Mutually-exclusive unary relations** label each domain element with the single property of being some $\sigma \in \Sigma$
- **Segments** in phonological applications



Unconventional String Models

- **Non-exclusive unary relations** allow distinct alphabetic symbols to share properties
- **Features** in phonological applications



(Strother-Garcia et al., 2016; Vu et al., 2018)